

Fortified Control-Plane Encapsulation with Session-Key Derivation for Secure IP Mesh Routing

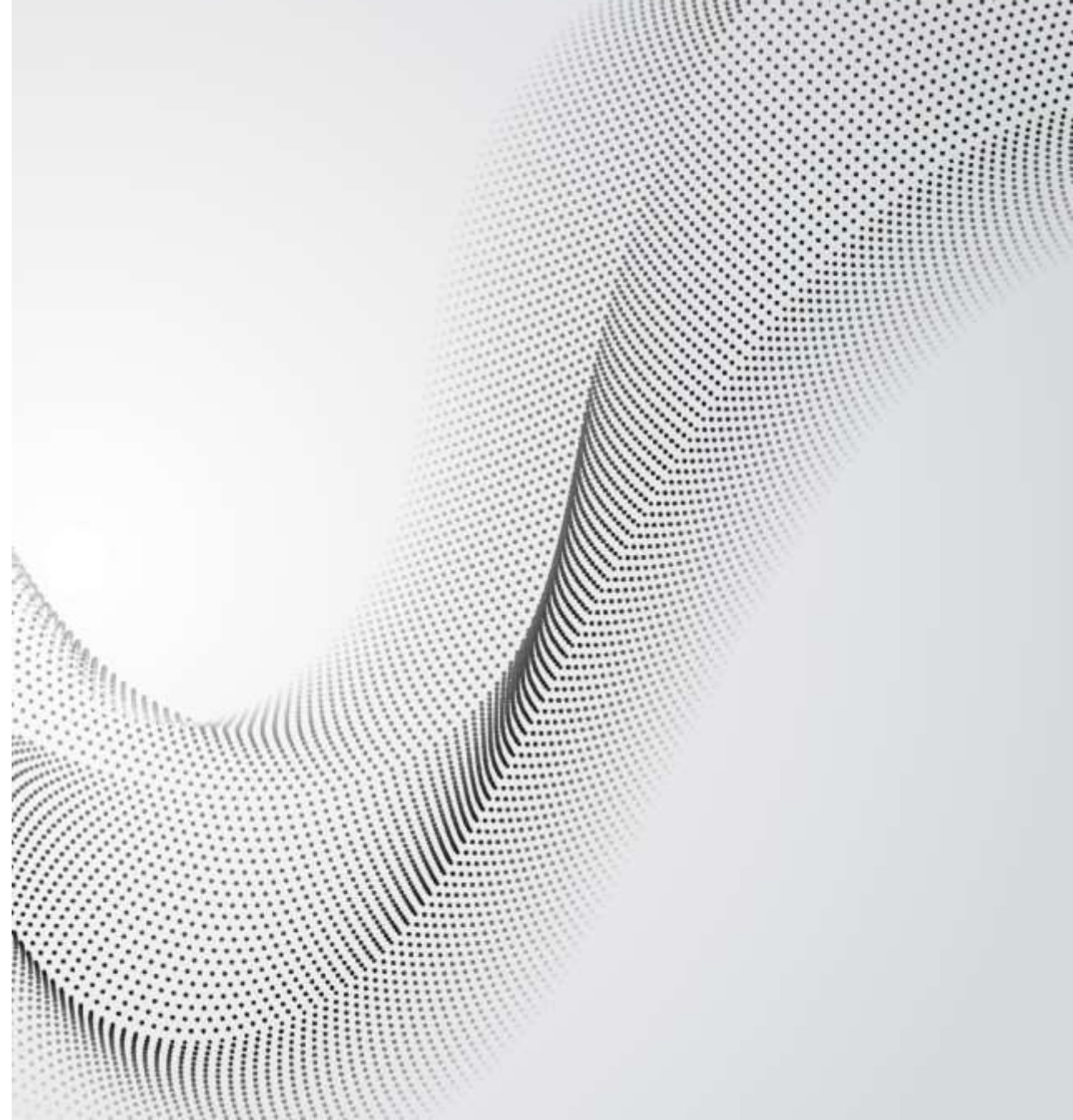
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Outline

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- ◆ OLSR & Quantum Cryptography Background
- ◆ PQC-Enhanced OLSR Design
- ◆ Implementation & Evaluation
- ◆ Discussion & Conclusions



Introduction & Motivation

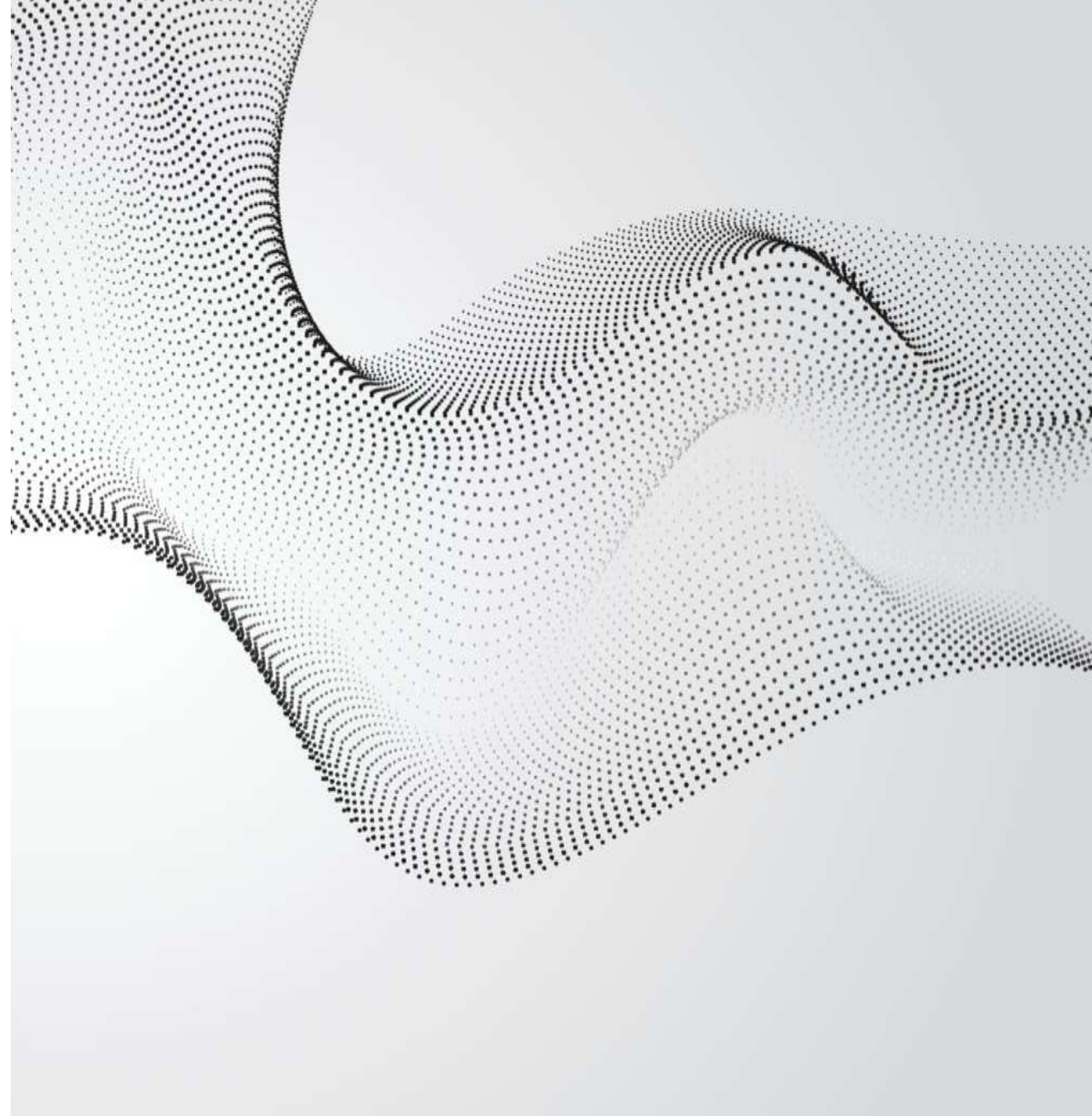
Mesh networks are widely deployed in community and tactical systems.

- Unencrypted OLSR control messages expose the network to eavesdropping and message forgery.
- Conventional security overlays rely on RSA/ECC, but these are vulnerable to quantum attacks (Shor's algorithm).
- The 'store now, decrypt later' threat means adversaries may record control traffic today to break it once quantum computers mature.

Issue	Impact
Unencrypted HELLO/TC	Topology leakage, spoofing
RSA/ECC reliance	Broken by quantum algorithms
Store-now-decrypt-later	Retroactive compromise of traffic

Background: OLSR & Mesh Networking

- ◆ Proactive link-state routing for MANETs via OLSR.
- ◆ HELLO messages discover neighbours and detect links.
- ◆ TC messages disseminate topology information via MPRs.
- ◆ Unencrypted control plane leaves mesh vulnerable.



Quantum Threat & Post-Quantum Cryptography

- Shor's algorithm breaks RSA/ECC, motivating quantum-safe cryptography.
- CRYSTALS-Kyber (Kyber512) is a lattice-based KEM selected by NIST for PQC.
- ChaCha20-Poly1305 AEAD offers 256-bit security and high performance without hardware acceleration.
- Together, Kyber512 and ChaCha20-Poly1305 provide confidentiality, integrity and authenticity against classical and quantum adversaries.

Primitive	Function
Kyber512 (KEM)	Session key establishment
ChaCha20-Poly1305	Encrypt & authenticate
HKDF	Derive symmetric key
TLVs 0x0F/0x10/0x11	Extend OLSR control fields

Architecture Overview

♦ Security Shim Layer

Acts as an interceptor between IP and OLSR logic. All control-plane packets pass through it for cryptographic processing.

♦ Inbound Processing

1. Detects presence of TLVs:

- 0x0F: Kyber Public Key
- 0x10: Kyber Ciphertext
- 0x11: AEAD Payload

2. Performs:

- Kyber512 key exchange (via liboqs)
- ChaCha20-Poly1305 decr. + tag verification

3. If valid → forward to OLSR core; else → discard.

♦ Outbound Processing

- ♦ Captures raw HELLO/TC payloads from OLSR core
- ♦ Encrypts using ChaCha20-Poly1305 with session key
- ♦ Appends AEAD TLV (0x11) to packet before IP transmission

♦ Key Material Management

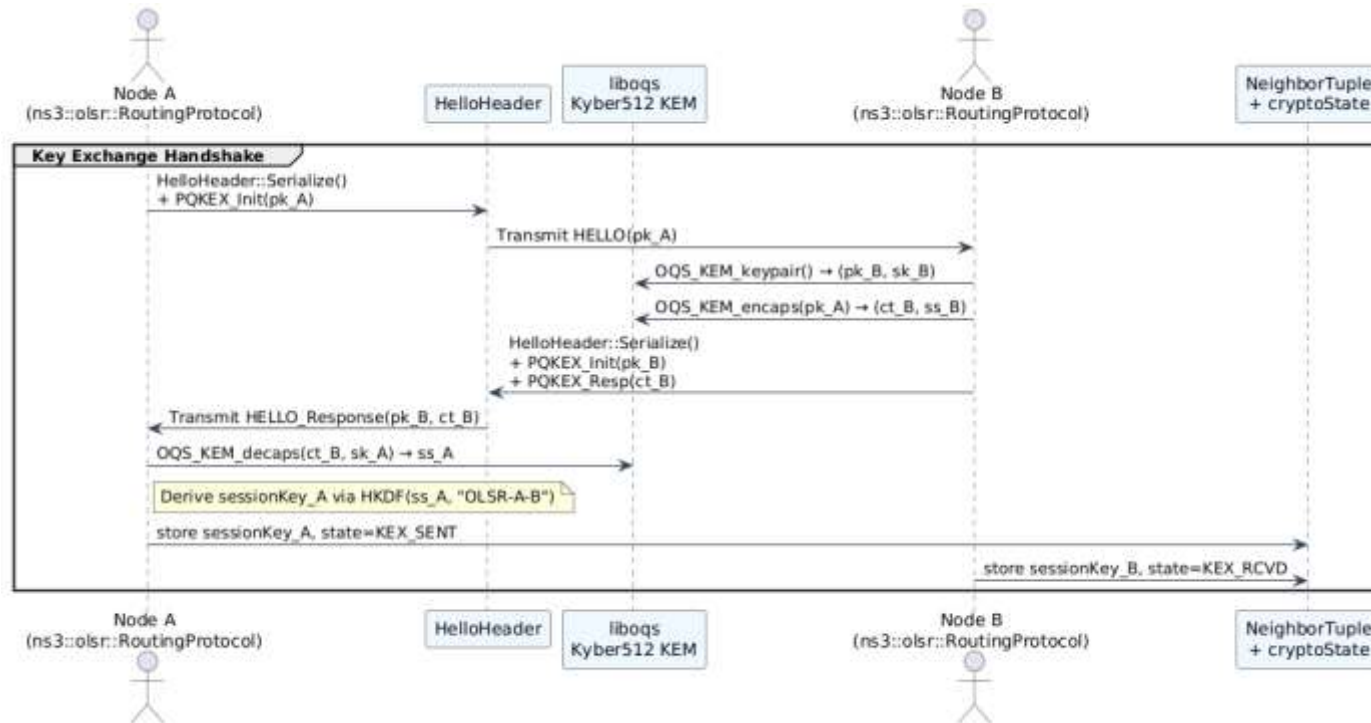
- ♦ Maintains per-neighbor state machine (e.g., INIT, KEX_RCVD, SECURE, BLACKLIST)
- ♦ Derives session keys via HKDF over Kyber shared secrets

♦ Compatibility

- ♦ Non-secure legacy nodes ignore unknown TLVs
- ♦ OLSR routing logic remains unchanged

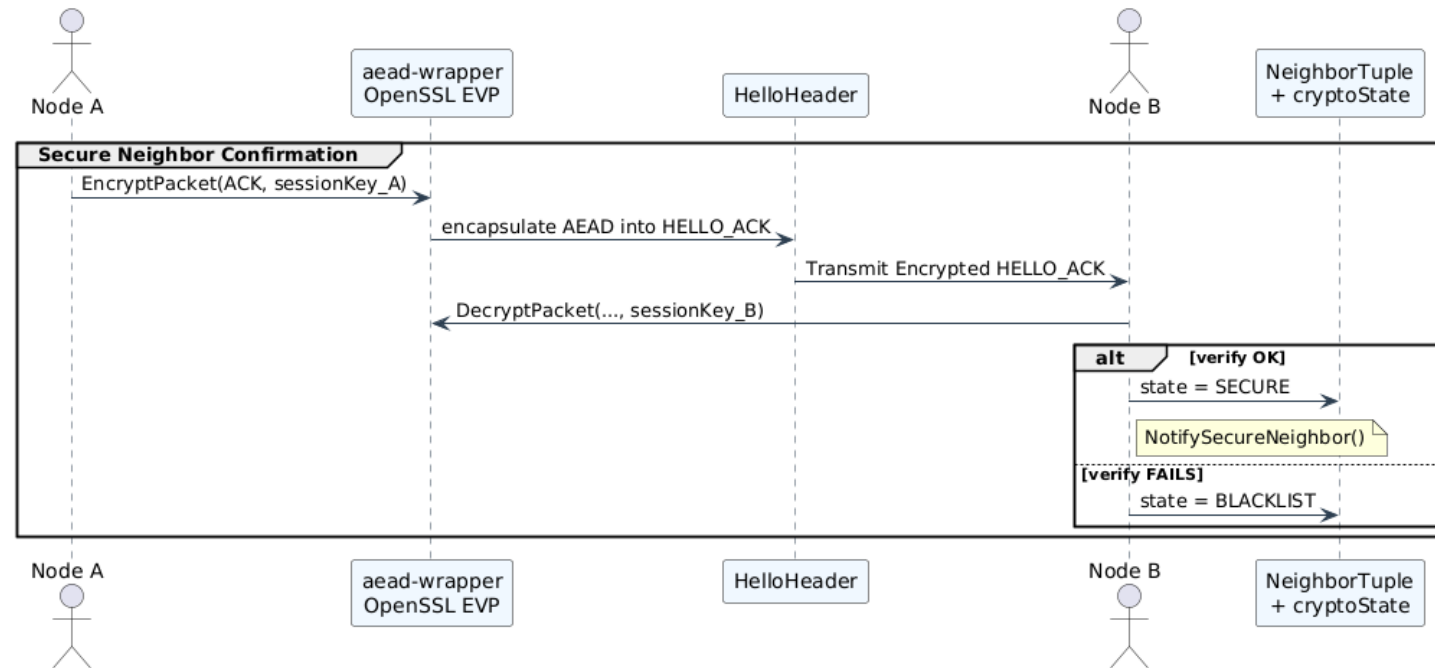
Element	Purpose
Security Shim	Between IP and OLSR, handles all crypto
Kyber512 (liboqs)	Key exchange via TLVs 0x0F, 0x10
ChaCha20-Poly1305	AEAD for HELLO/TC (TLV 0x11)
HELLO (Handshake)	Carries pubkey + ciphertext
HELLO/TC (Steady)	Encrypted with sesskey, AEAD-protected
Neighbor State	Tracks: INIT, KEX_RCVD, SECURE, BLACKLIST
HKDF	Derives sess key from shared secret
Failure Handling	Tag fail → drop packet, blacklist peer
OLSR Core	Unchanged; gets only verified payloads
Legacy Support	Unknown TLVs are safely ignored

Key Exchange Handshake



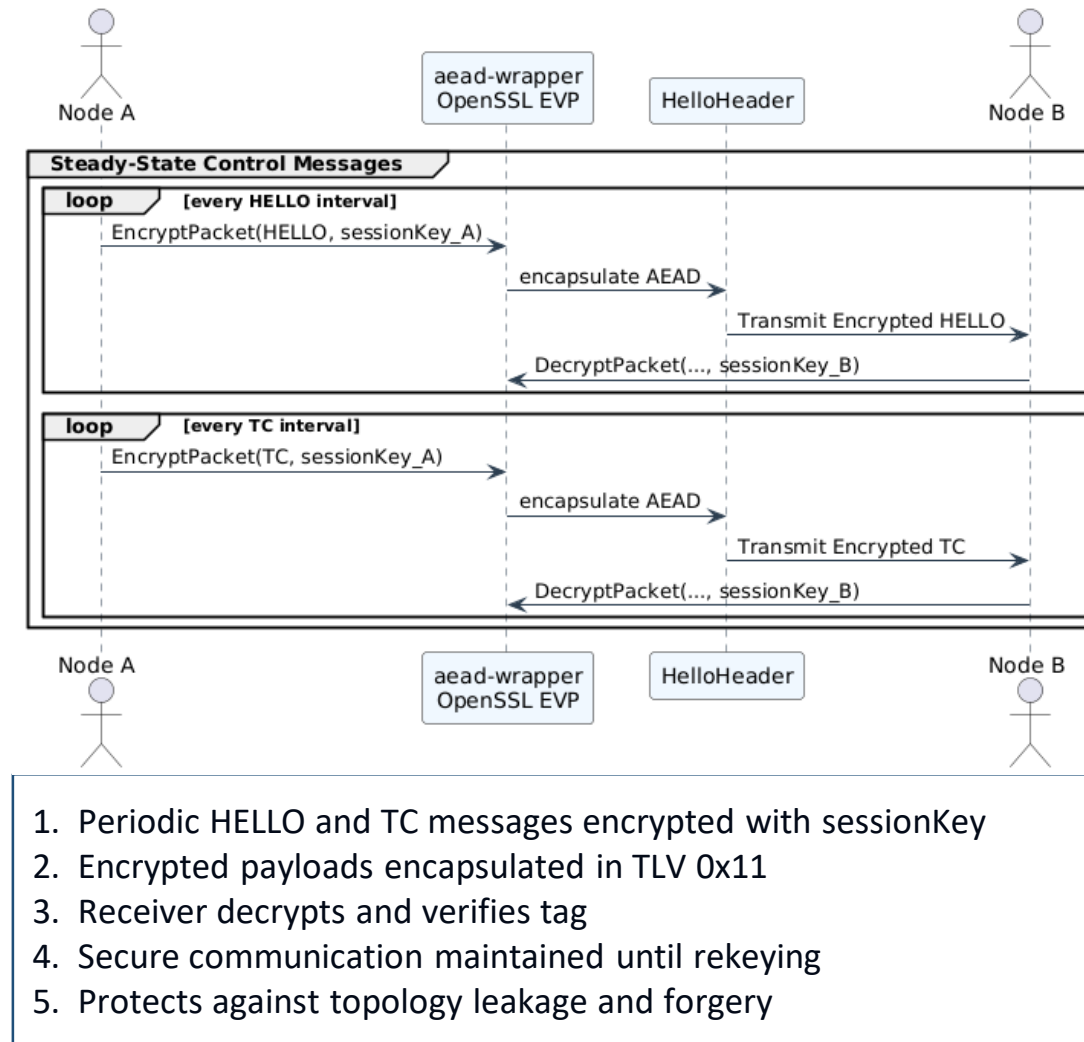
1. Node A generates a Kyber key pair (pk_A, sk_A) and sends pk_A.
2. Node B responds with its pub key pk_B and a KEM ciphertext ct_B derived from pk_A.
3. Both nodes derive a shared secret via HKDF and store a session key.
4. Node A sends an AEAD-protected ACK to confirm secure neighbour.

Secure Neighbor Confirmation



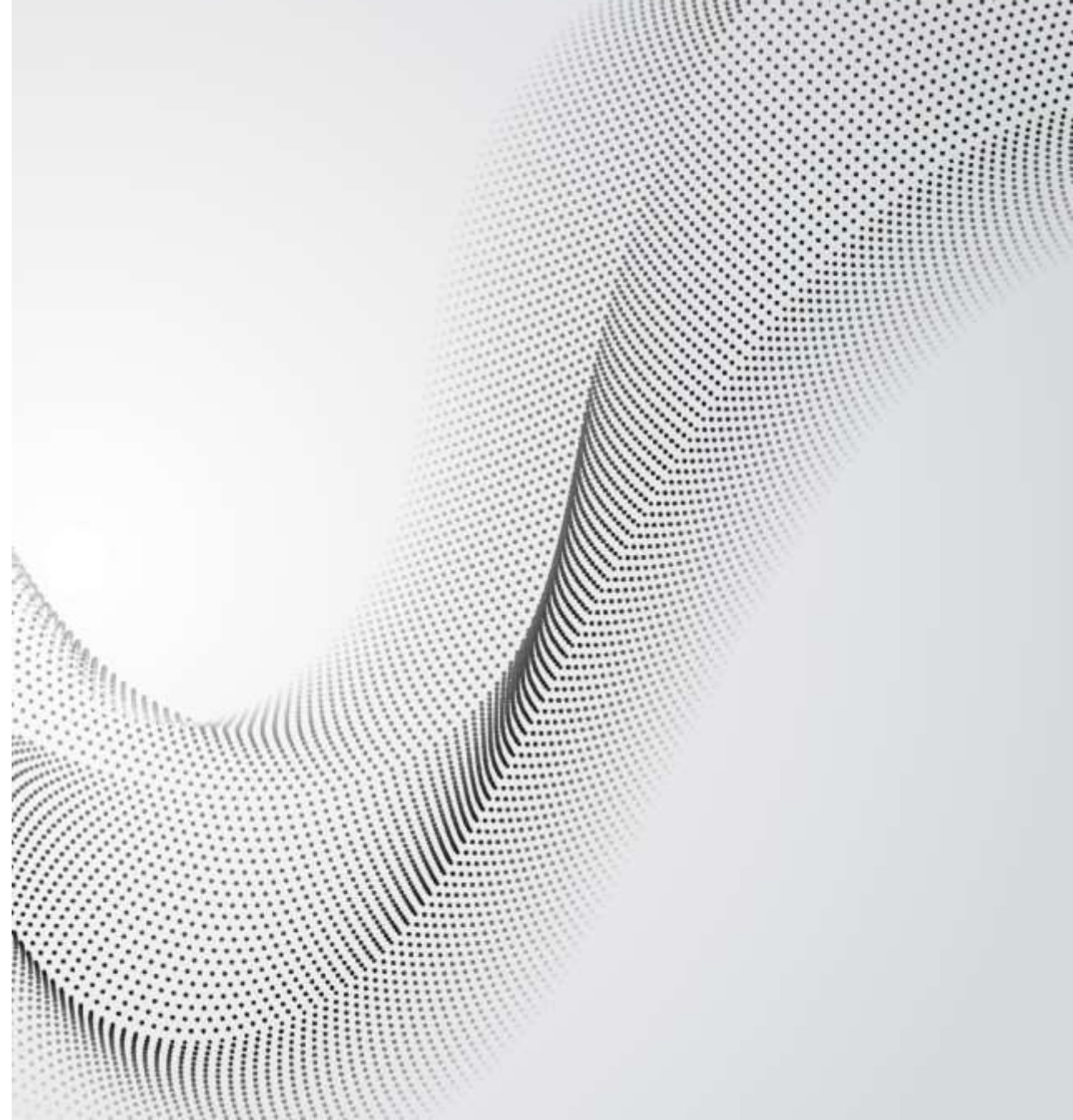
1. Once the key is established, every HELLO and TC msg is wrapped with ChaCha20-Poly1305.
2. An AEAD TLV (0x11) carries a 12-byte nonce, the ciphertext and a 16-byte auth tag.
3. Receivers decrypt and verify tags; invalid messages are dropped
4. The data plane remains unaffected because only control packets are encapsulated.

Steady-State Control Messages



Implementation & Evaluation

- ◆ NS3-based simulation framework
- ◆ OLSR HELLO and TC message re-structuring
- ◆ Evaluation in high- and low-mobility scenarios
- ◆ Dynamic overhead evaluation based on network size.



Evaluation Setup

Network & Mobility

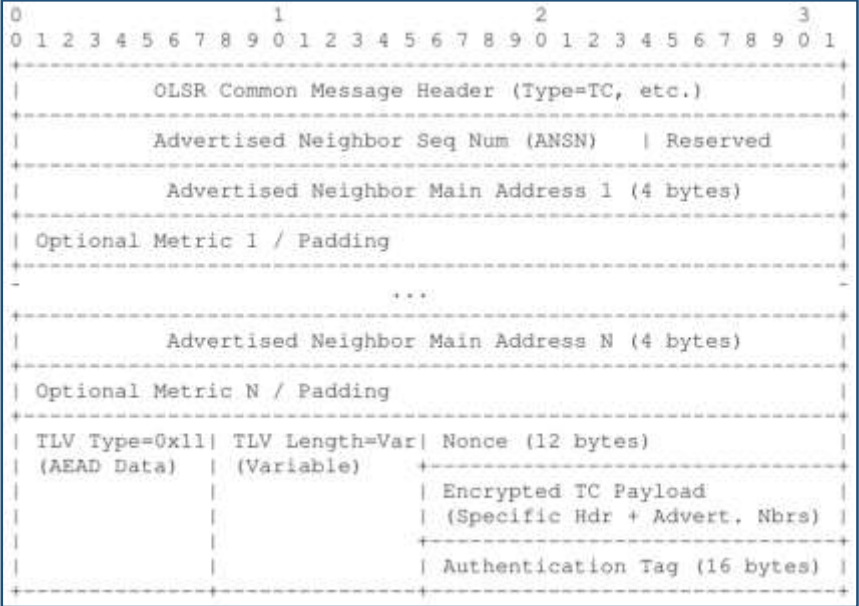
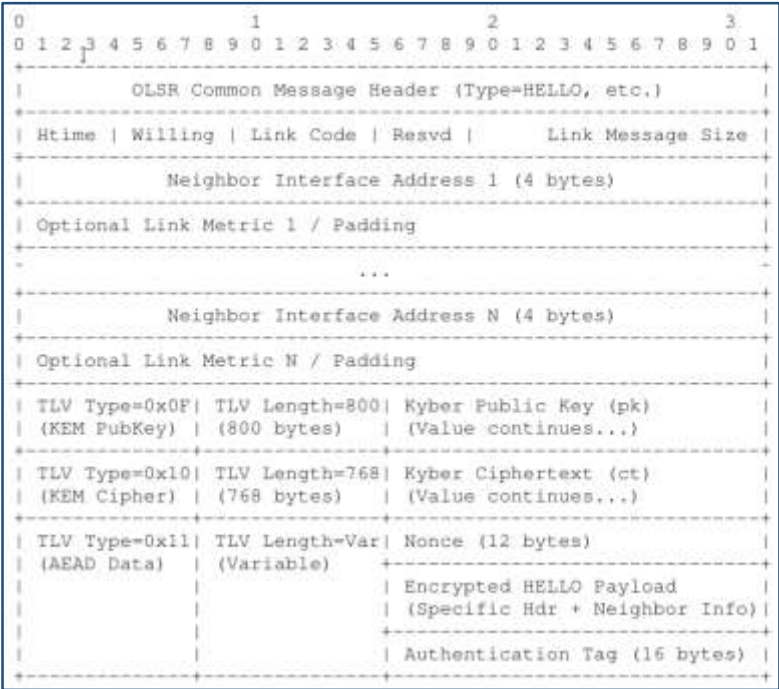
- Nodes placed in a 500×500 m area (static grid or Random Waypoint mobility).
- Mobility speeds: low mobility (1 m/s) and high mobility (10 m/s).
- Network sizes from 10 to 50 nodes.
- IEEE 802.11g WiFi with log-distance path loss.

Simulation & Metrics

- Instrumented control packets to measure per-node overhead (bits/s).
- Handshake time measured from KEX_RCVD to SECURE state.

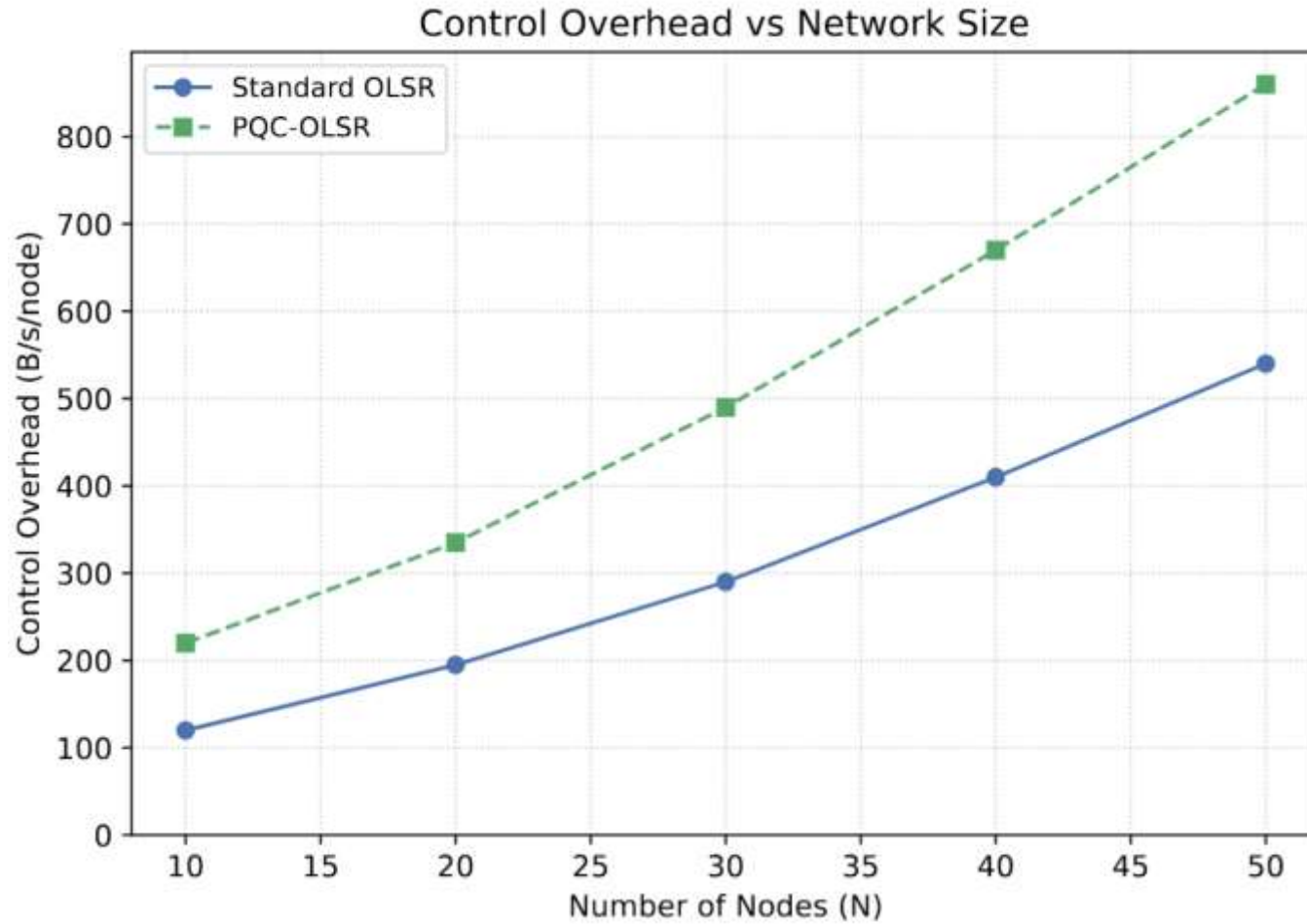
Implementation

- ♦ Integrated into the OLSR module as a security shim without altering core routing logic.
- ♦ Uses liboqs for Kyber512 KEM operations and OpenSSL EVP for ChaCha20-Poly1305.
- ♦ Public keys (800 B) and ciphertexts (768 B) exchanged via TLV 0x0F/0x10; derived keys drive AEAD.
- ♦ Handshake Phase (HELLO):
 - ♦ TLV 0x0F: Kyber512 Public Key (800 B)
 - ♦ TLV 0x10: Kyber512 Ciphertext (768 B)



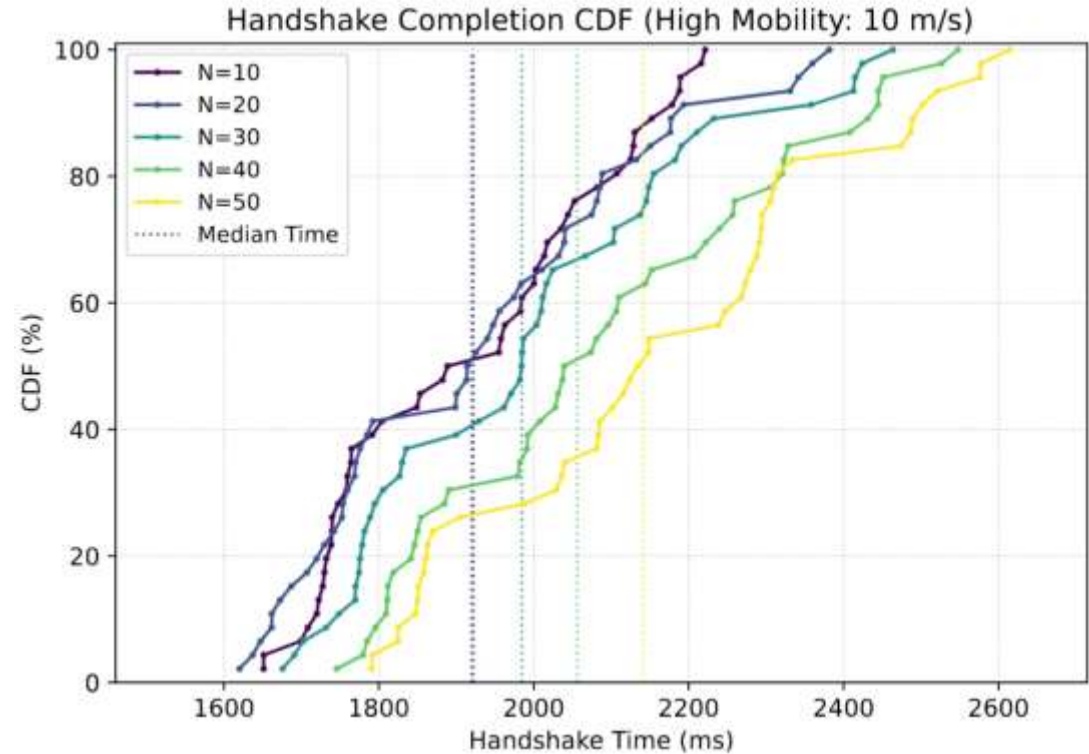
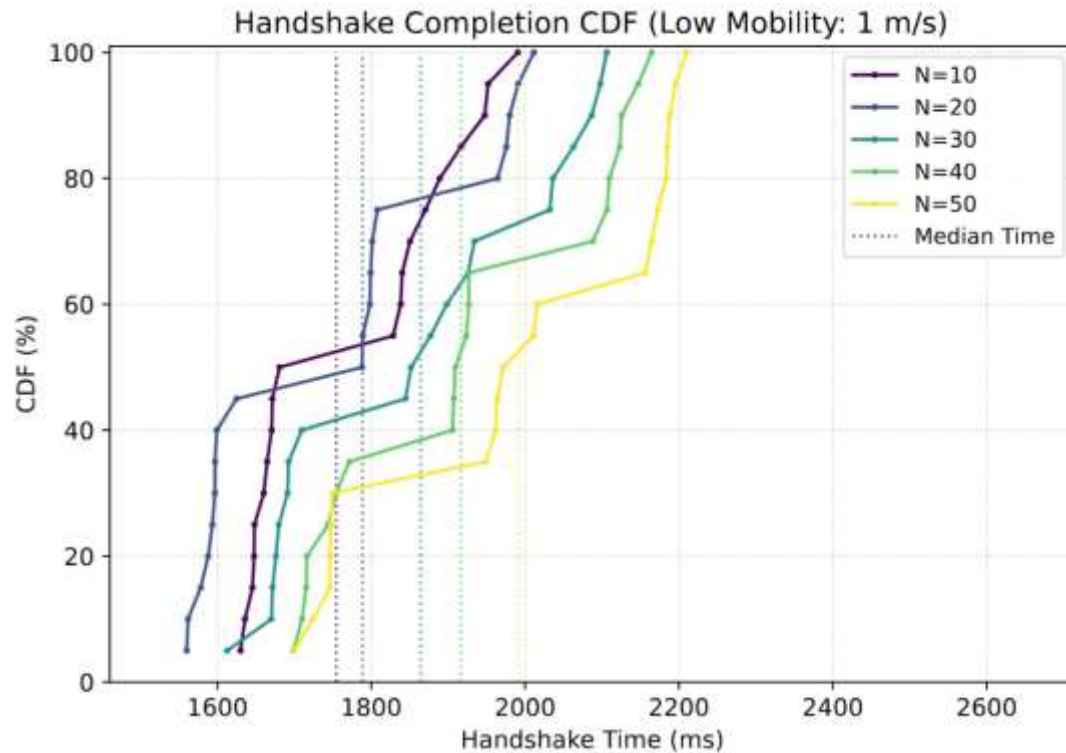
- ♦ AEAD overhead per message: 12 B nonce + 16 B authentication tag.
- ♦ Simulated 1 ms delay per PQC operation (Encapsulate/Decapsulate/Encrypt/Decrypt) to approximate computation cost.
- ♦ TLVs appended after core HELLO/TC headers without modifying base format.
- ♦ Backward Compatibility:
 - ♦ Legacy nodes skip unknown TLVs without breaking.
- ♦ Steady-State Phase (HELLO/TC):
 - ♦ TLV 0x11: AEAD payload → Nonce (12 B) | Ciphertext | Tag (16 B)

Results: Control Plane Overhead



- Baseline OLSR overhead increases with network size due to HELLO traffic and topology dissemination.
- PQC-OLSR overhead is significantly higher because each neighbour pair exchanges an 800 B public key and 768 B ciphertext during the handshake.
- Additional overhead comes from a 12 B nonce and 16 B tag appended to every secured HELLO/TC message.
- Despite the cost, the overhead scales linearly and provides confidentiality and integrity for the control plane.

Results: Handshake Completion Time



- In low mobility (1 m/s), handshake times cluster around 1.6–2.3 s; larger networks experience slightly longer delays.
- High mobility (10 m/s) shifts the distribution rightwards (1.8–2.6 s) and increases variance due to frequent neighbour changes.
- Even under dynamic conditions, handshake latencies remain within acceptable OLSR timescales.

Discussion

- ♦ Integrating PQC into the OLSR control plane is feasible but increases control overhead.
- ♦ Cryptographic delay has modest impact; handshake latency remains within protocol timescales.
- ♦ High mobility and larger networks amplify handshake variance, potentially delaying secure connectivity.
- ♦ Data plane performance is not affected (only control packets are encapsulated).
- ♦ Security gains (confidentiality, integrity, neighbour authentication) must be balanced against overhead in resource-constrained meshes.

Conclusion & Future Work

Conclusions:

- ✓ Presented a post-quantum secure extension of OLSR combining Kyber512 KEM and ChaCha20-Poly1305 AEAD via TLVs.
- ✓ Demonstrated practical implementation within ns-3 using standard cryptographic libraries.
- ✓ Evaluated control overhead and handshake latency across network sizes and mobility regimes.
- ✓ Security benefits outweigh overhead for many scenarios, providing confidentiality, integrity and authentication.

Future Extensions:

- Develop efficient rekeying and revocation mechanisms for dynamic networks.
- Design protocol-agnostic security layers to extend PQC protection beyond OLSR.
- Integrate Layer-3 control-plane security with data-plane protection for end-to-end security.

Thank you



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